

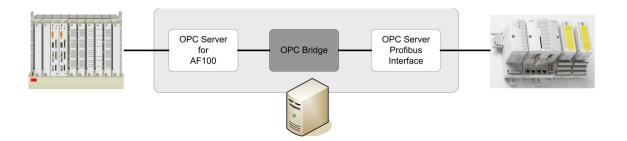
Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich Computer Engineering and Networks Laboratory



Master's Thesis

Industrial Networks Connecting Controllers via OPC

Performance – Availability – Reliability



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Abstract

In order to modernize their infrastructure and keep up with the state of the art, ABB Power Systems decided to replace the older controller AC450 with a new generation of controllers called AC800M. Just like its predecessor, its main task is to work as a sequencer in an otherwise mostly unchanging topology. Although the new controller AC800M provides modern communication features and a sophisticated application development system, it lacks of a communication interface compatible with the residing controllers AC160. A hardware approach addressing this problem is in development, but not available at this point of time. Thus the decision was made to realize the connection using OPC, a widely spread and open software communication interface standard with a high potential of reusability. In addition, it was aimed at gaining additional knowledge about the OPC interface, which is commonly used in industry.

In this thesis, we evaluate adequate hardware and software to realize this connection and we have programmed the controllers with applications to evaluate its performance and integrity. In addition, we are making considerations about redundancy that is vital in automation business in order to increase reliability and availability. We have shown that it is possible to interconnect controllers using OPC with satisfactory average performance results. Due to high maximum round trip times and high complexity when realizing redundancy, it is recommended to use such a system for testing purposes or non-critical operational applications, but not for critical systems. In this thesis we also identify and judge several alternative ways of connection.

Acknowledgements

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Chapter 1

Introduction

This chapter will provide a rough overview of the problem treated by this Master's Thesis. All technical devices and expressions will be explained more precisely in the next chapter. Please note that since this is a public thesis, it does not contain sensitive company-internal data.

1.1 ABB Power Systems

ABB Power Systems is one of the world's leading providers of infrastructure for controlling combined cycle power stations and waste-to-energy plants. Such a plant control infrastructure includes several hardware parts consisting of controllers, input/output-boards and communication devices as well as many software components to engineer, run, observe and analyze the power plant. A power plant control system has to satisfy a broad variety of different needs, from the efficient and reliable control of the turbines and associated supporting functions (such as lube oil) to easy configuration and operation as well as to sophisticated analysis functions addressing technical and economical aspects.

1.2 Problem Statement

Due to high investment costs, the technical management of power plants is a slowgoing business with long life-cycles. Thus, a considerable amount of hardware devices currently in use are tens of years old. For future applications within ABB Power Systems it will be necessary to connect a controller of the newest series used within ABB, Control IT AC800M, with an older controller of the type Advant Controller 160 (AC160). The problem is that these two controllers do not share a fast communication interface of similar type and therefore cannot communicate directly. The standard communication intended for AC160 is Advant Fieldbus 100 (AF100). However, AC800M can support a whole range of buses except for AF100. As a consequence, the communication must be implemented using some relaying technique.

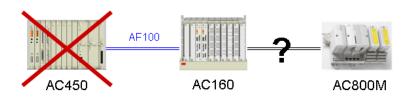


Figure 1.1: Basic problem: Communication between AC160 and AC800M

1.2.1 The Use of OPC

It was decided in advance to realize the relaying connection using OPC. This solution was chosen because OPC is an open standard and very common in process and automation industry. Furthermore, this solution offers a high potential to be used for similar problems, since a lot of devices support this specification. However, OPC is normally not used for fast controller-to-controller communication but for slower visualization and logging purposes and there is no performance data available for this kind of connection. The use of OPC is therefore both challenging as well as interesting to gain more insights and know-how.

It is also to mention that a hardware solution addressing our problem is not available yet. It is therefore necessary to have an alternative way using already available parts, also for testing purposes.

1.3 Goals

The goals of this Master's Thesis are stated as follows:

- Setup and evaluation of a test environment
- Setup of test systems
- Theoretical and practical evaluation of the test systems concerning performance, availability and reliability.
- Identification of improvements and different approaches
- Comparison with alternatives

As a starting point for the performance requirements, the current implementation was taken. The corresponding quantity and type of variables are displayed in Table 1.1 with 32-bit floating point values (floats) as analog in- and outputs and 1-bit boolean values as so-called status and command bits. In the current configuration with AC450 and AC160, all variables are written to the AF100 fieldbus with a cycle time of 256 milliseconds. Therefore we determined the minimum requirement for round-trip times from one controller to the other to exactly this time.

In agreement with the advisors, instead of elaborating the optional extension stated in the task description (Appendix C), we spent more time on trying out a

	AC160 to AC800M	AC800M to AC160
32-bit floats	600	32
1-bit booleans	2200	750

Table 1.1: Requirements for the quantity and type of variables [1]

second PROFIBUS approach and the theoretical derivation of a redundancy concept.

1.4 Structure

For the reader's convenience this Master's Thesis is structured thematically starting with an overview of components and terms (2) in the next chapter. The following chapters inform about the test system setup (3), the evaluations that were made (4) and finally the results (5). In a subsequent chapter the subject redundancy is treated (6) before the thesis comes to an end with the conclusion and outlook (7). Additional information as well as a CD-ROM containing more detailed data is located in the appendix of this thesis.

Chapter 2

Components and Terms

In this chapter, hardware and software parts as well as terms used for our test system and evaluations will be described. Some additional devices and programs concerning redundancy are introduced not until the chapter according. Information on the version numbers can be found in Appendix B.

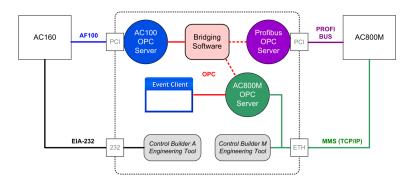


Figure 2.1: Schematic test system overview

2.1 Basic Terms

- **Performance**, in this thesis, refers to the capability of a communication component in means of speed and throughput.
- Availability is the term for the probability that a system will perform its specified functions when used under stated conditions. A common mathematical definition of operational availability is $A_o = MTBF/(MTBF + MDT)$, whereas MTBF is the "mean time between failure" and MDT the "mean down time" [2]. However, in this thesis, availability is used in a more general manner, since the basis for mathematical operations is not available.

- Reliability means the probability of a device remaining failure free during a specified time interval, e.g. the maintenance interval: $R = e^{\lambda t}$
- **Redundancy** is the implementation of extra components in addition to the ones needed for normal operation. Thus, redundancy normally increases reliability *and* availability.

2.2 OPC

OPC, originally short for "OLE for Process Control", is an open, standardized software communication interface specification launched in 1996 by a task force of different automation companies, later forming the OPC Foundation. As the former name indicates, OPC is an adaption of Microsoft's Object Linking and Embedding OLE¹ to the process control business, which used to be highly proprietary at that point of time. Thus it was almost impossible to efficiently combine products of different vendors. By providing so-called OPC servers with their devices, buses and software, vendors open their products to any OPC compliant client able to connect to the server for data exchange. Usually, an OPC server can handle several clients at once, while these clients—e.g. visualization or calculation applications—can connect to different servers in order to obtain their needed information.

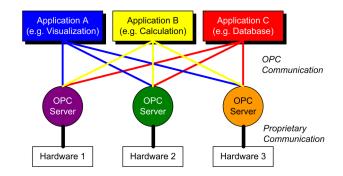


Figure 2.2: Classical OPC configuration

Over the years, the OPC Foundation has been adding eight additional specifications to the original one, therefore the name OPC was freed from its original meaning and is now used as an umbrella term [3]. Some important specifications are quickly explained in the following:

• **DA (Data Access)** is the original and most widely used standard of OPC. Its purpose is the cyclic polling of real time data, for example for visualization purposes.

¹http://msdn.microsoft.com/archive/default.asp?url=/archive/en-us/dnarolegen/ html/msdn_aboutole.asp

- HDA (Historical Data Access), in contrary, specifies the access to already stored data.
- AE (Alarms and Events) describes the non-cyclic, event-based exchange of alarms and events.
- **Data eXchange** is a specification from 2002 which regulates the direct communication between two OPC servers.

For this Master's Thesis it was made use both of the DA specification for the main purpose of communication as well as the AE specification in order to display and log round-trip times. Unfortunately, the promising Data eXchange specification is almost inexistent in practice and could therefore not be used in our thesis.

The underlying technique to exchange data is the component object model COM of Microsoft Windows, therefore OPC can only run on Windows operating systems [4]. A new generation of OPC specifications recently published is called OPC Unified Architecture (OPC UA) and is independent of COM, thus being able to run on more operating systems as well as embedded devices [5].

2.2.1 OPC Data Access

OPC DA is organized in the hierarchical structure server, group and item. Items correspond to variables and can be read and written. Furthermore, a quality and time stamp is provided with each of them. When reading items, the value usually comes from the OPC server's cache, which is updated periodically with the values of the device (or bus, component). However, it is usually possible to force a read directly from the device. Clients organize their items in groups, which for example share the same access method and update rate. Each OPC server has an unique name, some vendors even offer the operation of multiple servers for the same device. OPC DA provides different methods to access items, first of all synchronous and asynchronous read and write operations. More important to us, there is also a subscription mechanism, which is commonly used by modern clients in order to reduce communication. That is, the client group subscribes to the server which then "pushes" values towards the client only if they changed respectively exceed a pre-defined dead-band. The client can force an update of all these values by issuing a refresh call, which corresponds to an asynchronous read for all items of a group [6].

2.3 Programmable Logic Controllers

This section informs about the two controllers involved and about the controller that has to be replaced. Please notice that we use the term *controller* equivalent to programmable logic controller (PLC) throughout our Master's Thesis.

2.3.1 Advant Controller 160 (AC160)

The AC160 series was launched in 1997 to meet high speed requirements in turbine control. To this day its outstanding performance is needed for fast closed loop control (CLC). For our work, we were provided with a rack RF616 for the physical mounting of the controller parts. The rack also delivers power to each device and includes the BIOB Backplane Input/Output Bus which, among other tasks, processes the communication between the processor module and the communication interface. The tests in this Master's Thesis were done with processor modules of the type PM665 (containing a Motorola MPC8240 processor) and the AF100 communication interface CI631, both supporting redundancy [7]. To program the processor module, its built-in EIA-232 interface was connected to the engineering PC.



Figure 2.3: Advant Controller 160 (AC160)

2.3.2 Advant Controller 450 (AC450)

AC450 is the controller to be replaced by the new generation. Its processor module is built up around a Motorola 68040 microprocessor running at 25 MHz [8]. While its performance is quite limited, AC450 also lacks of modern communication interfaces and standards like Ethernet, which become more and more important.

2.3.3 Control IT AC800M

Control IT AC800M is the most current controller series used within all of ABB, introduced in 2001 [9]. It strikes with its small outline and offers modern communication interfaces. For our tests we were provided with a processor module PM864A which contains a Motorola MPC862 microprocessor running at 96 MHz [10]. In addition to two EIA-232 interfaces, the processor module of AC800M offers two built-in 10 Mbit/s standard Ethernet ports, used for instance to connect to the engineering computer.

The communication interface CI854A is a sophisticated device to establish PROFIBUS communication. The device even offers an own web interface which



Figure 2.4: Advant Controller 450 (AC450)

allows e.g. the surveillance of the cycle time. The device is directly coupled to the processor module via the Communication Extension Bus (CEX) [11]. At this point of time, CI854A can only act as a PROFIBUS master, but not as slave. As a consequence, communication partners must be slaves.



Figure 2.5: Control IT AC800M

2.4 Bus and Network Communication

This section describes the different communication approaches used in our system.

2.4.1 Advant Fieldbus 100 (AF100)

AF100 is a high performance fieldbus system with time synchronization and the most popular fieldbus used within ABB, common to almost all platforms. Although the name Advant Fieldbus 100 is ABB proprietary, it complies to IEC standard 61375-3 as well as US-standard IEEE 1473 and is widely used in railways under the name of MVB. AF100 works over twisted pair, coaxial and optical media at a rate of 1.5 Mbit/s. For our system we made use of the first method only. Bus master responsibility is shared by all communication interfaces with administrator

capabilities, that is, a bus master controls all communication for a certain time interval and then passes on the bus master token. However, in our setup only AC160 were given administrator capabilities, even though the PCI-card would support them as well.

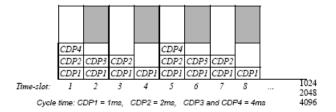


Figure 2.6: AF100 time-slot organization up to 2000 m bus length [12]

AF100 is a planned bus with a pre-determined scan table and thus meets realtime requirements. Process Data Transfer is managed through Cyclic Data Packets (CDPs). Each CDP is configured individually on the communication interface for a certain signal identity, cycle time, size and direction. Each broadcasted CDP has a unique signal identity, whereas receiving CDPs can have the same signal identity, provided they are situated in different communication interfaces. That is, multiple interfaces can receive the same CDP. The cycle time determines how often the data of the CDP is transferred on the bus. When a CDP is transferred on the Advant Fieldbus 100, the interval between consecutive transfers is always the same, the cycle time. Thus, process data transfer is deterministic, regardless of which other tasks the communication interfaces perform.

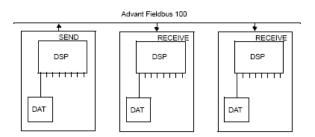


Figure 2.7: DataSet Peripheral and DAT Elements [12]

In AC160, one so-called DataSet Peripheral (DSP) can reference up to eight 32-bit values called Database Elements (DATs). Each DSP uses one CDP to be transported on the AF100, and can therefore be configured for example with an individual cycle time. In addition to cyclic message transfer, AF100 offers the possibilities to send event-based messages using the Service Data Protocol (SDP). This does not influence the cyclic data transfer at all. At a bus length up to 2000m it is therefore necessary to keep the cyclic bus load under 70% in order to have extra space for messages (depicted as gray fields in Figure 2.6) [12].

CDP size	DSP size	Transfer time
4 byte size	1 DAT element	0.088 ms
8 byte size	2 DAT elements	$0.108 \mathrm{ms}$
12-16 byte size	3-4 DAT elements	$0.156 \mathrm{ms}$
20-32 byte size	5-8 DAT elements	$0.252 \mathrm{ms}$

Table 2.1: CDP transfer time up to 2000 m bus length [12]

2.4.2 PROFIBUS DPV0

PROFIBUS (**PRO**cess **FI**eld **BUS**) is one of the world's most widely used fieldbuses and was launched in 1989 by a consortium of companies and institutions. Until today it was continuously enhanced and extended [13]. It is available in three variations:

- **PROFIBUS-FMS** (Field Message Specification) was the first PROFIBUS communications protocol. It offers sophisticated functionality but is no longer officially supported.
- **PROFIBUS-DP** (Decentralized Peripherals) has been designed for fast cyclic data exchange at field level and is today often referred to DPV0. This specification was extended with DPV1 for acyclic and alarm communication and DPV2 for slave-to-slave and isochronous communication as well as synchronized system time.
- **PROFIBUS-PA** (Process Automation) is used for the connection of devices to a superior process control system. It supports intrinsic safe transmission [14].

For this thesis, only the DPV0 protocol specification is relevant. Therefore, when writing PROFIBUS, we are always referring to this specification. DPV0 allows the cyclic exchange of data from a master (initiator) to slaves and vice versa at a high data rate up to 12 Mbit/s. Communication is organized in input and output modules of defined size (e.g. one word or four bytes) for every slave, which then are transmitted completely (non differential) once per cycle. The cycle time depends on the number of bytes and slaves as well as on timeout settings. The specified input and output data of one slave is limited to 244 byte each. While the physical connection can also be established using fiber optics, we used the more conventional EIA-485 electrical twisted pair connection.

PROFIBUS slaves come with a GSD-File ("Geraete-Stammdaten" in German) which is an electronic device specification. Most PROFIBUS master systems allow to directly import the GSD-Files of previously unknown slaves in order to include them in the bus planning. That makes it easy and comfortable to build PROFIBUS networks with slaves from different vendors.

2.4.3 MMS

The Manufacturing Message Specification MMS is an application level protocol developed during the eighties by General Motors. Its main goal was to unify communication to controllers independent of manufacturers. Other automobile, aerospace and PLC companies adopted the protocol and since 1990 MMS is standardized in ISO/IEC 9506 [15]. While being very general and open, MMS is reputed as heavy, complicated and costly. Nevertheless MMS was an important step in development and due to its independence of the underlying communication protocol it is still being used today. Furthermore it worked as a reference model and has influenced other protocols. Since 2002 the standard IEC 61850 for "Communication networks and systems in substations" based on TCP/IP over Ethernet defines the mapping onto MMS [16].

MMS is used as the main communication protocol for AC800M. On one hand, this is the communication between controller and the superior control and engineering system, on the other hand the communication between several controllers of this type as shown in Figure 2.8. One telegram can reach a size up to 1024 bytes including the header, containing a maximum of 159 integers or 136 floats [17].

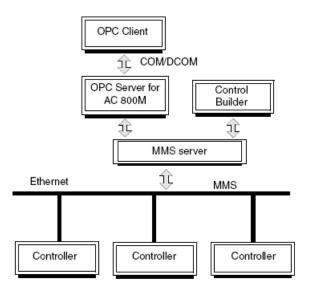


Figure 2.8: MMS communication for AC800M [18]

2.5 Personal Computer

For our tests we used a personal computer containing an Intel Pentium 4 processor at a speed of 3.2 GHz and 2 GB memory running a Microsoft Windows Server 2003 R2 operating system.

2.5.1 800xA and Development Environment

In order to setup, program and run ABB components, an *Industrial IT 800xA* environment was installed on the personal computer. Besides a lot of additional resources, this installation also encompasses the application development systems for both controller types.

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Figure 2.9: 800xA Engineering Workplace

The development environment to program the AC160 was Control Builder A (CBA) consisting of Application Builder, Function Chart Builder and Bus Configuration Builder. To program the AC800M controller, the engineering tool Control Builder M was used.

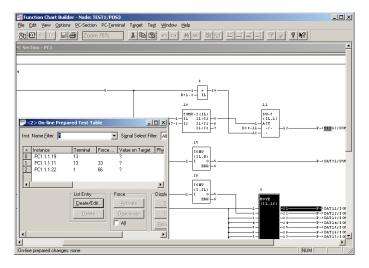


Figure 2.10: Function Chart Builder

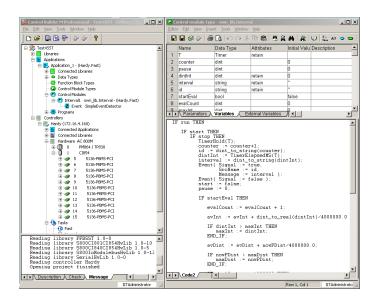


Figure 2.11: Control Builder M

2.5.2 AF100 Communication

To establish connection to the AF100 fieldbus, we inserted an ABB CI527 PCI card into the personal computer. The according AC100 OPC Server, which allows us to access the AF100 bus, was installed with the 800xA for AC100 software extension [19]. It is to mention that AC100 OPC Server allows access on bit-level, for example, an integer value is presented by the server both as integer value and split up in 32 boolean values.

2.5.3 MMS Communication

An Intel Ethernet PCI card allowed the communication with the AC800M via MMS on TCP/IP. The according *AC800M OPC Server* is part of the 800xA installation. All communication over this port is performed via the Manufacturing Message Specification (MMS) protocol running over TCP/IP, utilized for example by the engineering tool to program the controller. The same connection can also be used for controller to controller communication when having several MSS-ready devices. Furthermore, the AC800M OPC Server communicates with the controller via the same protocol and infrastructure, making available all variables by default [17].

2.5.4 Beckhoff PROFIBUS Communication

For the first PROFIBUS connection we used the FC3102 PCI card from Beckhoff. This card was chosen due to its flexibility: It provides two ports in one PCI card which can be freely adjusted either as master, slave or passive bus monitor [20]. We ran Beckhoff's TwinCat software: *TwinCat System Manager* to configure the card, *TwinCat IO* as driver and *TwinCat OPC Server*. Beckhoff inserts an abstraction layer on top of TwinCat IO which operates with the ADS (Automation Device Specification) protocol [21].



Figure 2.12: Personal computer backside

2.5.5 Woodhead PROFIBUS Communication

The Woodhead Electronics SST-PBMS-PCI card is the only card to be found that provides multi-slave functionality. That is, even though the card only has one physical SUB-D9 interface, it can emulate up to 125 PROFIBUS slave interfaces [22]. This allows not only the simulation of large PROFIBUS outlines in one PC, but—and more important to us—also makes it easy to bypass the 244 byte limitation specified by PROFIBUS DP.



Figure 2.13: Woodhead Electronic SST-PBMS-PCI card

We configured the card with its configuration tool SST Profibus Configuration. Furthermore we used the according OPC Server for PBMS Card which was configured with Multi Slave OPC Configuration Utility.

2.5.6 OPC Bridging Software

The programs we used to interconnect two OPC servers were Matrikon's *OPC Data Manager* (ODM) [23] and Kepware's *LinkMaster* [24]. These programs called OPC routers or OPC bridges are able to read data from one server and write it to another. Both programs are similar in configuration and operation. The functionality includes the definition of groups and update rates, input/output pairs, dead-bands and quality checks. LinkMaster even allows to write one input value to more than one output variables and to perform mathematical operations in between. To make bulk configuration easier (e.g. with Excel), both programs allow to import and export the configuration from and to comma separated values (CSV) files.

	Users Tools	er\TEST1.lmf] (Demo Expires 01:58:00) Help			
) 🛩 🔛 🐚	i 💱 😭 🗠	3 🖻 🖻 🗙 🔔 📓			
E 🚇 Local Mach E 💼 Remote M E ݵ Custom Re	achine	E 🚽 Local Machine E 😤 Remote Machine E 📽 Custom Remote			
🛃 simob	Link Name	Input	Outputs	Data Type	Raw Value
🛃 simor	somib0	Local Machine\ABB.AC100.1\SOMIB_0.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
a somib	Somib1	Local Machine\ABB.AC100.1\SOMIB_1.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	1416
somir	Somib10	Local Machine\ABB.AC100.1\SOMIB_10.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib11	Local Machine\ABB.AC100.1\SOMIB_11.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib12	Local Machine\ABB.AC100.1\SOMIB_12.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib13	Local Machine\ABB.AC100.1\SOMIB_13.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib14	Local Machine\ABB.AC100.1\SOMIB_14.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib15	Local Machine\ABB.AC100.1\SOMIB_15.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib16	Local Machine\ABB.AC100.1\SOMIB_16.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib17	Local Machine\ABB.AC100.1\SOMIB_17.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib18	Local Machine\ABB.AC100.1\SOMIB_18.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib19	Local Machine\ABB.AC100.1\SOMIB_19.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib2	Local Machine\ABB.AC100.1\SOMIB_2.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib20	Local Machine\ABB.AC100.1\SOMIB_20.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib21	Local Machine\ABB.AC100.1\SOMIB_21.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib22	Local Machine\ABB.AC100.1\SOMIB_22.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib23	Local Machine\ABB.AC100.1\SOMIB_23.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib24	Local Machine\ABB.AC100.1\SOMIB_24.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib25	Local Machine\ABB.AC100.1\SOMIB_25.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib26	Local Machine\ABB.AC100.1\SOMIB_26.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib27	Local Machine\ABB.AC100.1\SOMIB_27.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
	Somib28	Local Machine\ABB.AC100.1\SOMIB_28.VALUE	Local Machine\BECKHOFF.TwinCATOpcServerDA\InOu	DWord	2220510
🛃 Link Groups	l la î				

Figure 2.14: Kepware LinkMaster GUI

We ran both bridging programs with a fully functional, time-limited testing license provided for free by its vendors for the duration of our thesis.

2.5.7 Helper Programs

For setup and testing, a range of other software was used on the engineering/test system computer. The most important programs are shortly specified here:

- MatrikonOPC Explorer is a freeware OPC client allowing to connect to any compliant OPC server and displaying the value of chosen tags. It also supports writing of variables and preserving settings. Furthermore, it allows measuring the maximum update rate of the OPC servers it is connected to.
- Office 2003 of Microsoft was used for day to day work and configuration tasks. Especially *Excel* was helpful for variable definition in AC800M and for bulk configuring the bridging software using CSV files. Furthermore, with the

Bulk Data Manager plugin, Excel allows convenient bulk setup of PROFIBUS devices and its variables.

- **Process Explorer** of Microsoft's Sysinternals is a freeware program to monitor system processes and their use of resources.
- **Synergy** was used to operate two computers and screens with only one keyboard and mouse.
- Irfanview helped with the caption of screenshots.

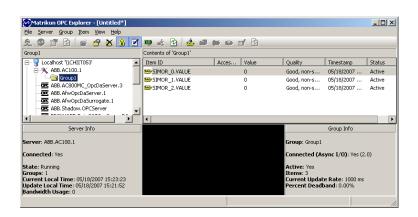


Figure 2.15: MatrikonOPC Explorer

Chapter 3

Test System

This chapter describes the test system we used for our evaluations.



Figure 3.1: Picture of the test system

3.1 Overview

The main goal is to connect a controller of the type AC160 with a controller of the type AC800M which do not share a common communication interface. As this happens using OPC, the intermediate relay is implemented using a Windows computer. While there is only one possibility to connect AC160, we tried out different approaches for linking AC800M.

3.2 AC160 Side

The AC160 is attached to the personal computer via the CI631 using the Advant Fieldbus 100 and a PCI card CI527. Due to limited communication facilities, this is the only reasonable choice for this task. Since both communication interfaces inherently support line redundancy (see Chapter 6), two media were used for this

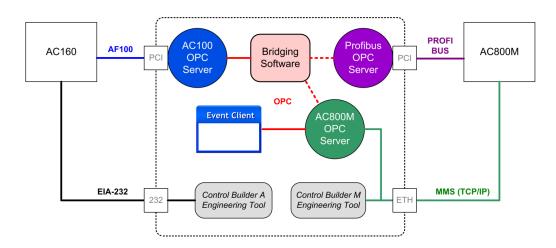


Figure 3.2: Schematic test system overview

connection. To program the processor module either a separate EIA-232 connection or the AF100 fieldbus itself was used. The according AC100 OPC Server is part of the 800xA for AC100 extension but can also be installed stand-alone [25].

All DSPs sent from AC160 are available in the AC100 OPC Server if not set to be filtered out by the configuration tool. It is also possible by default to write to these variables, however, if not configured otherwise the values written on the OPC Server are not sent using planned DSPs but using event-based Service Data Protocol (SDP) messages. As a consequence, all communication towards AC160 becomes very slow and is not even guaranteed to arrive. It is therefore necessary to implement *receiving* DSPs in the AC160 and to configure the AC100 OPC Server to send DSPs. The latter is done using the Bus Configuration Builder [26]. In the properties dialog of the OPC station it is possible to specify the DSPs that have to be sent [27]. Using this method, data transmission towards the AC160 becomes as fast as reading from it.

Station Properties	×					
Station Type : AF100 OP0 Station Number : 15	:					
This AF100 OPC station will distribute acknowledge signals for DAT-based objects to other AdvaSoft or OPC-server nodes with the following station numbers:						
This AF100 OPC station creates sending DSPs DSPs with one of the following values at DSP: 20 22						
OK	Cancel					

Figure 3.3: Configuration of sending DSPs in Bus Configuration Builder

3.3 AC800M Side

For all of our evaluations we connected the AC800M with our personal computer using Ethernet. This is firstly necessary for programming and managing the controller. Secondly, we displayed the measured results using OPC AE and therefore needed this connection in order to let the AC800M OPC Server communicate with the controller. Last but not least we tested this connection for our purposes as described in the following subsection.

3.3.1 Communication via MMS

MMS using TCP/IP over Ethernet is the standard communication between computer and controller, and since with AC800M OPC Server a ready-to-use connection is provided, we included this approach in our evaluations. However, the performance of this connection depends directly on the CPU load of the AC800M as well as on the Ethernet network traffic [28].

3.3.2 Communication via PROFIBUS

Due to the missing real-time behavior of a standard Ethernet link, it was planned right from scratch to test this approach against a PROFIBUS connection. With the CI854A, a sophisticated PROFIBUS communication interface for AC800M exists, supporting 12 Mbit/s baud rate not depending on the CPU of the processor module. This reasons and the fact that it enjoys a high acceptance in entire Europe made PROFIBUS to our first choice of all fieldbus systems. Another promising bus, FOUNDATION FIELDBUS High Speed Ethernet (FF HSE) was dismissed after having recognized that the according communication interface CI860 has a very poor throughput of only 500 32-bit values per second [17].

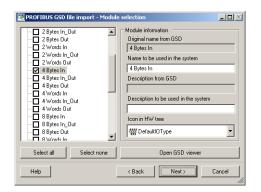


Figure 3.4: Module selection in Device Import Wizard

A typical use of PROFIBUS is to connect several slave devices such as sensors with one or few masters (for example a controller) on one single bus (up to 126 devices). Often the amount of data per slave is very small, since it only represents an input/output state or a measured value. Therefore it is rather exotic to use PROFIBUS as a point-to-point communication with a big amount of data as we did in our case. Since the first test showed that the Beckhoff PROFIBUS solution has performance problems, it was decided to test a second PCI card/OPC server product from Woodhead Electronics. Both PROFIBUS cards were configured in a way that one module consists of four 8-bit bytes to transport integers or two 16-bit words to transport floats. These mappings were realized using the Device Import Wizard in CBM and the PCI card's configuration tools [29].

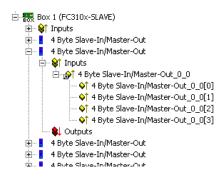


Figure 3.5: Cofiguration of PROFIBUS modules in TwinCat System Manager

3.4 Bridging Software

A core part of our system is the OPC bridging software establishing the connection between two OPC servers. Two different products were tested: OPC Data Manager from Matrikon and LinkMaster from Kepware. Both products were evaluated due to their specifications, bulk configuration possibilities (CSV files import/export) and availability for testing.

There are different possibilities for connections in OPC Data Access. The standard method implemented in both programs is *subscription*, that is, the bridging software subscribes as a client to the the source OPC server A to read from. The source server then *pushes* changing values to the client at a fixed speed (10 milliseconds is the maximum for LinkMaster, 1 ms for ODM) which writes them to the destination server. As a consequence, the load and therefore the performance of the bridging software depends on the number of changes.

Since OPC DA is organized in subscription groups, the signal pairs in the bridging software were combined in groups of 10 to 32 signals with identical update rates of 10 milliseconds. Both programs also support the (faster) processing of arrays instead of single values. However, we could not take advantage of this solution since the OPC servers do not support arrays.

Link Group Properties	×
General	
Identification	
Name: simob1	
Description:	
Settings	
Server <u>u</u> pdate rate: 10 🛖 milliseconds	
Client I/O refresh rate: 0 👘 milliseconds	
Enable link transfers	
OK Cancel Help	

Figure 3.6: LinkMaster link group properties

3.5 Resulting Test Systems

This section shows more detailed models of the resulting three test system setups following the considerations above. The AC160 side (illustrated on the left hand side) stays the same for all three setups. The models visualize one main challenge of our task, namely the big quantity of different interfaces and each component running with an own unsynchronized cycle time. While the models are not complete, they include all parts that are considered relevant in relation to our time domain milliseconds. It is to be mentioned that the impact of BIOB and CEX on the performance is considered to be very small compared to the other parts.

3.5.1 System 1: Beckhoff PROFIBUS

Figure 3.7 shows the schematic setup for the Beckhoff PROFIBUS approach. It is to mention that Beckhoff uses their proprietary ADS protocol to communicate between the OPC server and the card driver, which was therefore outlined in this figure [20]. The maximum update rate of TwinCat OPC Server is 2 milliseconds, the cycle time of the variable i/o task is 1 ms.

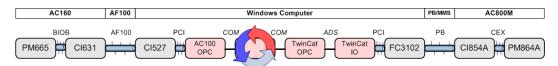


Figure 3.7: Schematic setup of Beckhoff PROFIBUS approach

3.5.2 System 2: MMS

The MMS approach is depicted in Figure 3.8. In contrary to the PROFIBUS solution there is no need for an additional communication device connected to

CEX bus as the Ethernet port is located in the processor module itself, using its resources. Unfortunately, the maximum update rate of the AC800M OPC Server is limited to 50 milliseconds, which is a comparatively high value compared to the cycle time of other system parts.

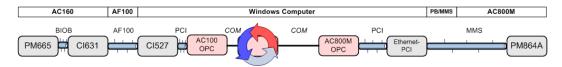


Figure 3.8: Schematic setup of MMS approach

3.5.3 System 3: Woodhead PROFIBUS

In contrary to Beckhoff's solution, Woodhead implements a more direct access path for the OPC server. Its maximum update rate is 1 millisecond. The GSD-file of the SST-PBMS-PCI card had to be slightly modified in order to be importable by Control Builder M. Furthermore, there exists a trick how to simplify bulk setup of OPC tag names. More information on the latter as well as a working GSD-file is enclosed with the CD-ROM in the appendix.

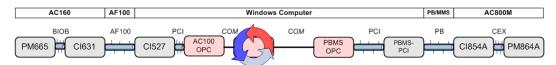


Figure 3.9: Schematic setup of Woodhead PROFIBUS approach

Chapter 4

Evaluation

In this chapter we describe the evaluations performed on the test system. In order to implement the different tests, we needed to program both controllers with appropriate code, which presumed the know-how for two quite different application development systems and engineering tools (see [30], [31], [32] and [33]).

4.1 Test Systems

As described in the preceding chapter, we evaluated three different test system variants. While the AC160 side physically stayed the same for all tests, the AC800 connection was tested with the following approaches:

- PROFIBUS connection with Beckhoff's card and server
- Standard MMS connection using an Ethernet card
- PROFIBUS connection with Woodhead's card and server

Due to reasons explained in the next chapter, all main tests were performed with Kepware's LinkMaster OPC bridging software.

4.2 Basic Idea

The basic idea behind all tests is the definition of a complementary set of receiving and sending variables to each of the two controllers. During each controller program cycle all receiving variables are read from the bus/network and all sending variables are written. On the other side of the bus/network, the OPC server makes the variables available. The OPC bridging software, acting as an OPC client to the servers, links the two corresponding variables by reading the values from one server and writing them to the other.

AC100 OPC Server	Bridging Software	AC800M OPC Server
Sending Integers		Receiving Integers
Receiving Integers	<	Sending Integers
Sending Floats		Receiving Floats
Receiving Floats		Sending Floats

Figure 4.1: Variable bridging in normal mode

4.2.1 Short Circuit

To explore the behavior of one side or to isolate problems, we also performed test with a "short circuit" configuration. That is, the bridging software is not configured to link the variables from one OPC server to the other, but right back to other variables on the same server.

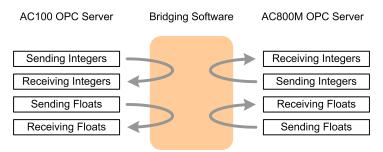


Figure 4.2: Variable bridging in "short-circuited" mode

4.3 Variables

This section informs about the type and quantity of variables used to test the connection.

4.3.1 Data Types

According to the requirements, only 32-bit floating point and boolean values have to be processed. Since the booleans are packed into integer values, the problem modifies to the transport of 32-bit floating point and 32-bit integer values.

4.3.2 Packaging of Booleans

The are various reasons to pack booleans into integer values instead of processing them one by one:

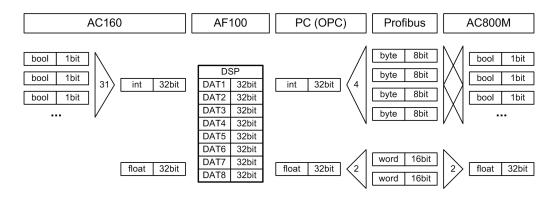


Figure 4.3: Data types in different sections with PROFIBUS

- The smallest possible PROFIBUS module size provided by both cards is one byte. If processing booleans one by one, 7 bit would get "lost" for every value, which reduces performance.
- In MMS telegrams, one single boolean requires 3 bytes space whereas one integer containing up to 32 booleans uses only 6 bytes [17].
- A smaller number of variables improves the performance of the OPC connections.
- One can save a lot of configuration effort for the bridging software.
- AC800M supports *transparent* packaging of the booleans into PROFIBUS modules, that is, it is not even necessary to manually pack them into integers.
- AC160 offers a program code element to pack/unpack booleans into/from integers.

Since the size of one variable in an AF100 DSP is limited to 32 bit, we intended to pack this amount of booleans into an integer value. However, the packaging program code element offered in AC160 treats the last bit as a sign flag not compatible to the transparent mapping of the AC800M. Instead of implementing additional program code we therefore decided to set this last bit aside and pack only 31 bits into one integer value.

4.3.3 Quantity

In order to learn more about the behavior of the concerned parts, we performed six main tests with different amounts of variables:

• The first four Tests 0 to 3 are intended to get an overview on the behavior of the configurations. They are designed such that they can run with all test system variations, that is, the amount of data does not exceed 488 bytes in one direction (due to the limitation of two PROFIBUS slaves with Beckhoff's

solution). Test 0 is carried out with only one integer value going forth and back. Test 2 contains twice and Test 3 about three times the quantity of variables as in Test 1, allowing to observe linear tendencies.

- Test 4 represents the *real* requirements for the quantity of variables and data types according to Table 1.1. Thus, the results of Test 4 are the most important ones for statements regarding the feasibility of the setup. In contrary to Tests 1 to 3, the amount of floating point values transmitted from AC160 to AC800M is much higher. As a consequence, it cannot be performed with Beckhoff's PROFIBUS solution.
- Test 5 is the duplication of Test 4, trying to determine the limits of the system.

	Test 0	Test 1	Test 2	Test 3	Test 4	Test 5
AC800M to AC160	1	20	40	60	56	112
AC160 to AC800M	1	42	84	122	670	1340
Total	2	62	124	182	726	1452

Table 4.1 :	Number	of 32-bit	signals	for	each	test

4.4 Measurements

The following methods were used to evaluate our test system.

4.4.1 Round-Trip Time

The most relevant evaluation we performed is the round trip time (RTT) measurement. For this reason, a boolean value (packed in an integer) is set true and sent from the AC800M controller to the AC160. At the same time, a stopwatch is started. Inside the AC160 the value is moved to another boolean and sent back. As soon as this value returns to AC800M, the stopwatch is stopped and the RTT displayed via generated event that allows subsequent interpretation for 500 entries. This number is the minimum sample quantity for our results discussed in the next chapter. The actual communication slowed down by the cycle time of AC160 and the resolution of the results is limited by the cycle time of AC800M. The results can therefore be looked at as worst-case values, while the exact time used for communication only is considered to be somewhat lower in most cases.

An important thing to know is that all the other values, which are transmitted with every cycle but are not directly relevant for measuring the RTT, changed *every* controller cycle. While this does not matter for the fieldbus communication it has an influence on the OPC communication. The common subscription mode for the OPC bridging software works differential, that is, values are only transmitted when they change, which saves CPU load. In order to test the worst case scenario—which is our task—it is therefore necessary to transmit a new value with every cycle.

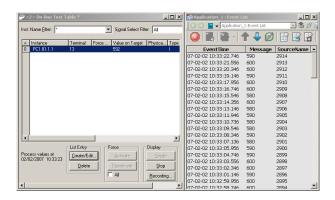


Figure 4.4: Measuring RTTs on AC160 and AC800M

A problem that showed up was the cycle time of the AC160 program code and the cycle time of DSPs transmitted on AF100. For Tests 0 to 3, both values are constantly 8 milliseconds, which make these tests very comparable. Since the amount of values increases a lot for Tests 4 and 5, the performance of both the CPU and the bus are not sufficient anymore and need higher cycle times. However, higher cycle times on AC160 and AF100 make the RTT slower and results are not directly comparable to the first tests due to changed conditions. For better understanding it was therefore decided to test with different configurations and cycle times depicted in Table 4.2. The trade off is to have lower (faster) cycle times in exchange for OPC server load: While with configuration A it is a worst-case test, the RTTs get very high and are not directly comparable anymore. Configuration C makes also Test 4 and 5 directly comparable to the others, but it is not a worst-case test anymore.

For better understanding the following definitions we introduce two terms: For each configuration there is only *one receiving* and *one sending* DSP relevant for measuring the round-trip times on AF100, containing the variables to time the connection. We call them the two *measuring* DSPs while all other DSPs are called *load* DSPs.

- **Configuration A:** The cycle time for *all* DSPs on AF100 is identical and complies with a safe choice for the AF100 bus load (see Appendix D). Since for Tests 4 and 5 the program code runs faster than the bus, there is a new value sent in each bus cycle.
- **Configuration B:** The cycle time for the measuring DSPs is set to same value as the program code cycle time. All load DSPs are sent as in configuration A. For Tests 4 and 5 this means that load values change only every fourth time in comparison with the measuring value.
- **Configuration C:** It is of course not possible to set the DSP cycle time faster than the program code. Thus it would never be possible to reduce it to 8 ms for Tests 4 and 5. We therefore installed a second processor module, working

at 8 ms cycle time and only processing the measuring DSPs while the load DSPs were still processed in the first module. For Tests 4 this means that load values change only every eighth time compared to the measuring value, for Test 5 it is even only every 16th time.

	Test 0	Test 1	Test 2	Test 3	Test 4	Test 5
Configuration A	8	8	8	8	64	128
Configuration B	8	8	8	8	16	32
Configuration C	8	8	8	8	8	8

Table 4.2: Cycle times of measuring DSPs

The 8 ms cycle time for AC160 program code and AF100 for Tests 0 to 3 was chosen since it fits for all these tests and still is in the range of other component's cycle times. Of course it would have been possible to simply increase the cycle times for Tests 0 to 3 in order to have the same preconditions as in Tests 4 or 5. But since it is our task to evaluate the performance of the *communication* and not the controller we decided not to falsify results by doing so.

4.4.2 Integrity

The intention of this measurement is to prove that every linked value is communicated correctly (not just the ones relevant for measuring the RTT). For this reason, all float variables were set to exactly the same value, being increased by one at every cycle and sent from one controller to the other. On this receiver side, the range between the maximum and minimum value are identified. If the range is small, integrity is good. If the range is wide, this shows that some variables do not change at the same time and therefore are not transmitted correctly.

4.4.3 Availability

To get an idea about the availability of such a system, one configuration was running for one week (seven times 24 hours). The goal of this evaluation was to find out about the capability to run the system for a longer period or if it fails after some hours. Furthermore, the two preceding measurements are included with recording the average and maximum RTT as well as average and maximum integrity check range. Due to time limitation this availability long-time test was only performed using the Woodhead PROFIBUS approach and a realistic variable quantity according to Test 4 Configuration B. In order not to impair the measurement and falsify results, we abandoned a more detailed logging.

While numerical information on availability of industrial devices it based upon extensive tests and long-time information, this was not within the scope of our thesis. A mathematical approach cannot be performed either, due to the lack of numerical information of the single parts in the personal computer.

Chapter 5

Results

In this chapter we will present and discuss the results gained from our evalutations.

5.1 Bridging Software

We started the evaluation using Matrikon's OPC Data Manager. Unfortunately, the behavior of the GUI was sometimes unsatisfying, particularly due to its long delays. Additionally, the communication stopped at times for unknown reasons. The equivalent product LinkMaster from Kepware offers similar possibilities in configuration but did not show these effects. Furthermore, the performance of LinkMaster could be determined as being vastly better than the one of ODM: For Test 3 with Beckhoff PROFIBUS the median RTT with LinkMaster is 510 milliseconds, whereas with ODM and the same preconditions the RTT becomes more than ten times slower. As a consequence, all main tests were performed using Kepware's LinkMaster software. For some configurations we enabled a device refresh rate of 10 ms for only one group in each direction, which prevented the program from unwanted stopping.

5.2 Round-Trip Times

This section shows the results for the round-trip time tests, grouped by test system variations. In Table 5.2 we show the expected maximum RTTs based on the addition of cycle times and update rates.

Test 1	Test 2	Test 3	Test 4	Test 5
$2.7 \mathrm{ms}$	$3.2 \mathrm{ms}$	$3.5 \mathrm{ms}$	$7 \mathrm{ms}$	$10 \mathrm{ms}$

Table 5.1: Rounded PROFIBUS cycle time measurements

For the PROFIBUS cycle times we entered the maximum values measured with CI854A shown in Table 5.1, Test 3 for the Beckhoff setup and Test 5 for Woodhead

	Beckhoff	MMS	Woodhead
AC160 code	8	8	8
AF100 DSPs	8	8	8
AC100 OPC	10	10	10
LinkMaster	10	10	10
OPC server	2	50	1
i/o task	1		
PROFIBUS	3.5		10
AC800M code	10	10	10
Total*2	105	192	114

approach. The configuration of the PROFIBUS parameters was done automatically by AC800M with "actual value based" method for all test we performed.

Table 5.2: Expectation of maximum RTTs Configuration C

5.2.1 System 1: Beckhoff PROFIBUS

Evaluations were performed for different amounts of signals as defined in Table 4.1. Starting with the Beckhoff PROFIBUS solution, we could observe that the RTT increases linearly to the amount of signals. While test results with one single signal (Test 0) were promising, the following tests were already unacceptable, although defining only a fraction of the required amount of variables.

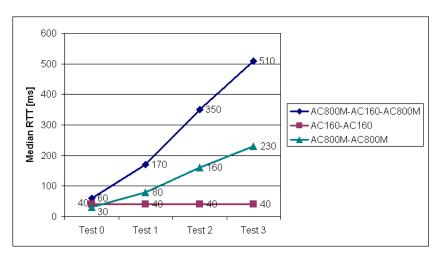


Figure 5.1: First test results with Beckhoff PROFIBUS

After short-circuiting the connections it was quickly determined that the bottleneck is to be found in the writing to the AC800M. After inspecting the LinkMaster software and PROFIBUS configuration the bottleneck was isolated as the OPC server to PCI card communication. Interrogating Beckhoff support confirmed this suspicion: Since it is much more common to read from OPC servers, the reading process was optimized by transporting several values in block mode from the PCI card to the OPC server, using Beckhoff's own ADS protocol. In contrary, values written to the OPC server are transmitted one by one [34]. This explains our results, and after putting into operation two more clone OPC servers, which is the maximum number, the RTT expectedly was more or less three times lower. However, with 190 ms median RTT for Test 3 and added configuration effort, this is still not satisfactory. It is to add that Beckhoff offers a nice and sophisticated configuration tool, which also allows the swapping of words and bytes. This function is used when processing float values which are interpreted different in FC3102 PCI card than in CI854A.

After these findings it was inevitable to look for another solution, also keeping thoughts in mind that it would be necessary—due to the 244 byte limitation—to put *six* FC3102 PCI cards into a computer in order to have enough capacity for all required values, or alternatively to have *eleven* DP/DP-couplers at work¹. The idea to "multiplex" data over the PROFIBUS in order to defuse byte limitation was discarded as impracticable.

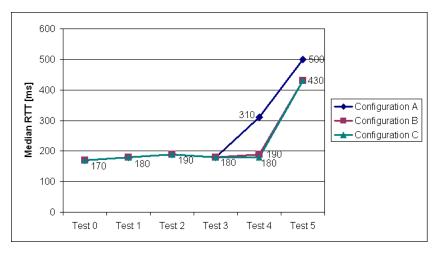


Figure 5.2: Test results with MMS

5.2.2 System 2: MMS

Before looking at the results, it is to mention that, compared with the PROFIBUS solution, the MMS approach suffers from vast disadvantages: Firstly, the maximum update rate of the AC800M OPC Server amounts to 50 ms and is therefore a big obstacle. Since MMS cannot perform much better anyway (we estimate around 40 transactions per second [28]) this rate is justified, though. Secondly, the performance depends on the processor module CPU and network traffic. While the

 $^{^{1}}$ DP/DP-couplers are hardware relaying devices able to connect two PROFIBUS DP masters, acting as a slave device on each side and copying data back and forth.

median RTTs with Configurations B and C meet our requirements, the results could not convince us entirely. As suspected, the standard deviation of the measurements is enormous. For different configurations of Test 4 it varied in the range from 40.9 to 62.7 ms.

It is not due to MMS that the results for Test 5 show such bad results. This time it is the fault of our personal computer which is not able to handle this amount of variables anymore. That is, the CPU load reaches 100% and therefore slowing down the LinkMaster process.

5.2.3 System 3: Woodhead PROFIBUS

Since the SST-PBMS-PCI card supports the definition of several virtual slaves in one card, the 244 byte limitation is not a problem any more. Thus, also Tests 4 and 5 could be performed with this one single card. This card and its server show a good performance in both directions even with a big quantity of variables. Of course, the CPU overload problem with Test 5 shows up here as well.

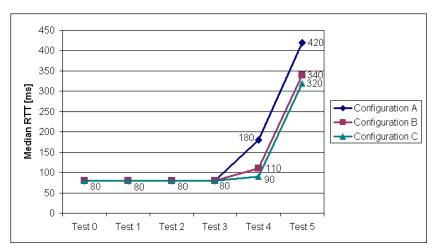


Figure 5.3: Test results with Woodhead PROFIBUS

While the performance is satisfactory, another problem could not be solved well: The word order in CI854A is interpreted differently than in the SST-PBMS-PCI card. More precisely, 32-bit floats are split in two 16-bit words on PROFIBUS and are thus sent in a certain order (see Figure 4.3). They are re-assembled in the opposite interface just the other way round, such that floating point values are always interpreted wrong. According to Woodhead Electronics and ABB, this alleged small problem cannot be solved changing the configuration of either part. At the present day it is therefore necessary to implement additional program code in AC800M to manually perform word swaps for each float. It is to add that we experienced the configuration tools offered by Woodhead as being of less convenience compared with Beckhoff's solution. Especially the naming of tags in big quantities is troublesome, but fortunately there could be found some workarounds [35].

5.3 Integrity

The results of our integrity test for selected configurations are very satisfying. The maximum range of the received variables never exceeded the value to be expected according to the maximum RTT. As an example, for Test 4 Configuration B a maximum range of 6 in one direction is considered a normal delay due to transmission time and not indicating an integrity problem of the bridging software or different parts. More detailed, since we measured a minimum RTT of 60 ms and a maximum RTT of 290 ms for this specific setup (results contained on CD-ROM), the difference of the fastest possible way and the worst path in one direction is considered about (290ms - 60ms)/2 = 115ms. Within this time-frame, the program code (cycle time 16 ms) is executed 7.2 times. That is, if two variables arrive at the same time, but one took the fastest and one the slowest path, a maximum difference of 7 is to be expected. However, while the integrity check showed that things run correctly for different variants and configurations in a short-term view, its results of the long-time test covered by the next section might even be of more importance.

5.4 Availability

The long-time test lasted for exactly one week and ran without interruption, therefore proving the basic long-term functionality of the system. Within this time, a quantity of 2,283,900 samples was taken and analyzed.

	Average	Maximum
Round-trip time	$119.3 \mathrm{ms}$	$2759~\mathrm{ms}$
Integrity range	6.03	168.0

Table 5.3: Results of long-time test

The contents of Table 5.3 show that the results are satisfactory in average. With an expected maximum RTT of 146 ms according Configuration B^2 , the average RTT value is inside the expected range. The average integrity range fits to this value and affirms the correct transfer of values also in a long term. However, the maximum RTT of almost three seconds and the consequential integrity range is not what we had expected from our system. Most likely this maximum RTT are due to temporary CPU overloads caused by concurrent processes on the personal computer, or because of an OPC server failure. If an OPC server fails, LinkMaster tries to restart it after some time, however, this time was adjusted to 10 seconds for our evaluations.

In contrary to the RTT test above, all floating point signals in AC800M had to be processed one by one in order to switch the word order. This was necessary for the integrity evaluation, otherwise the signals would have been interpreted wrong,

 $^{^2\}mathrm{Replacing}$ the first two values with 16 ms in Table 5.2

causing a worthless range. This extra program code slowed down the cycle time such that it was not possible anymore to keep 10 ms, which of course also influences the RTT results. With an average cycle time of 12 ms, the impact is still not devastating, though. A maximum cycle time of 43 ms was displayed at the end of the test.

5.5 Summary

Figure 5.4 shows the combined results for the different variants. For better comparability, only Configuration C is displayed, such that the relevant DSP cycle time is 8 milliseconds for all tests (see Table 4.2). It shows clearly, that the test system variation with Woodhead PROFIBUS returns the best round-trip times, while the Beckhoff approach is worthless for our task. Moreover, its ability of emulating several slaves saves the effort of inserting multiple PCI cards. Beside all these advantages, the Woodhead PROFIBUS approach causes the word swap problem, which cannot be solved without adding program code at this time.

Whereas MMS also meets our requirements, it is about two times slower than the Woodhead approach. Furthermore, the standard deviation of the RTTs measured is vastly greater with MMS than with PROFIBUS. This is remarkable, since our test were made with a direct Ethernet link (crossed cable). As the performance of MMS depends on network traffic and AC800M's CPU load, it is to be expected that communication using MMS will get slower and even less deterministic in a real environment.

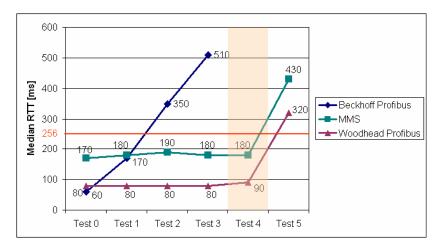


Figure 5.4: Test results summary (configuration C)

Although on the first spot Test 5 seems to be useless, this is not the case at all. Firstly, it proved that the system basically works even with this amount of variables. Secondly, while Test 4 is the most relevant evaluation concerning the real requirements, Test 5 was intended to look for the limits of the current system,

and so they were found. However, using a more powerful CPU would allow setting up a system with this quantity of variables.

Chapter 6

Redundancy

This chapter will provide theoretical considerations on redundancy. Since the Woodhead PROFIBUS solution was the most promising approach, our considerations are based on this test system variation and the amount of signals according to Test 4, which resembles the real requirements.

6.1 Terms and Concepts

Please notice that when talking about redundant components in this context, we mean *duplication* if not stated otherwise. The term *redundancy* does not define the number of redundant components in general.

6.1.1 Levels of Redundancy

There is a large variety within the levels of redundancy. While it is possible to double only the most important or critical components, one can also make a whole system redundant. This usually allows the concurrent outage of several different parts of a system (e.g. a processor module and a bus line) without interrupting the whole system. An *outage* can be an unplanned failure of a component but also a planned maintenance action. Security relevant parts like protection systems are often even implemented three times or more.

It is also important to know that redundancy can be implemented on different levels in communication. Most redundancy hardware devices or software programs that are able to switch between different connections work on a protocol level, that is, they recognize errors in the communication protocol e.g. if a device fails or the connection is broken. They are not able to identify errors in the transported data, though. In contrary, redundancy logic on application level checks the values/contents of the transported data, and therefore also detects errors if the protocol itself runs correctly. It is also possible to combine these approaches, e.g. if an application level program is able to force a switchover in a physical switch. A typical specification for the maximum switchover time in the turbine control business is 20 ms.

6.1.2 Master-, Slave- and Line-Redundancy

In bus systems there are three aspects of redundancy which can be combined arbitrarily. Since the outage of a non-redundant bus master in a classical master/slave bus system like PROFIBUS interrupts any communication, it is very common to implement bus master redundancy. In contrary, slave redundancy can be implemented depending on the importance of a specific slave. For real master or slave redundancy it is of course necessary to implement two completely independent communication stacks [36]. In practice this is normally done by using two identical communication interfaces. The two interfaces communicate directly or via a third component (e.g. the processor module) to be activated or deactivated and exchange configuration data. Line redundancy means the multiple presences of physical media. A redundant master/slave bus system is showed in Figure 6.3. ABB offers a device called RLM01 to connect single-line PROFIBUS slaves to a line redundant layout, allowing reducing the non-redundant part to a short distance. Since the device needs to copy data, it causes a delay, but this is very small compared to our cycle times (about 1.5 μ s at 12 Mbit/s).



Figure 6.1: ABB's PROFIBUS Redundancy Link Module RLM01

6.1.3 Transparency

For the different components of a redundant system it makes a difference whether the redundant pairs behave in a transparent manner or not. If they behave transparently, the other parts recognize a redundant pair as one single device and will not have to do anything; at the most there will be a more or less drastic pause due to the switchover. If redundant components behave not transparently, it is usually needed that nearby components are intelligent enough to check which device is working correctly.

For redundant components which behave transparently it is inevitable to have a redundancy communication (RedCom), which can be a direct link or be managed by a superior component. This connection is on one hand needed to determine which component is the active one and which is in stand-by mode. On the other hand, also status information has to be communicated. If a component fails it can for example be very important that the redundant part is able to recover the most recent values in order to take over the communication.

6.1.4 Stand-by modes

There are different modes of redundancy. Sometimes, the redundant parts run in parallel (load balancing) but can take the whole load in case the second system stops working, as for example with water pumps or internet web servers. This kind of redundancy is sometimes also referred to as *hot standby with load balancing*. With line redundant bus systems, the sending component usually writes on both lines whereas the receiving part chooses the line to read from.

In contrary there exist modes where only one part of a redundant pair can be active, while the other part is in a standby mode ready to take over. Dependent on the level of stand-by we talk about cold, warm or hot stand-by. *Cold stand-by* usually means that the redundant system first has to start up and gather necessary information before it is able to take over the communication, causing a comparatively long delay. *Hot stand-by* in contrary means that the redundant part is all the time up and running and even knows status information of the active system, always ready to take over the shortest possible amount of time. *Warm stand-by* at last is something in between.

In bus systems, where normally the different bus subscribers are addressed, the "flying redundancy" concept as in the CI854A can be used. That is, the active communication interface and the backup interface always have the same address, independent of which is the active one and which the backup. As a consequence, the active and the passive interface have to switch addresses if a switchover takes place. The advantage of this implementation is, that the backup slave has a valid address and therefore can be reached as well, to receive status information, for example. Moreover, such a switchover is transparent for the other bus subscribers. However, if there is an error during the switchover, it can happen that both redundant components have the same address, the effects of which would be devastating.

6.1.5 Power Supply Redundancy

A lot of controllers and automation components are prepared to be connected to two different power sources to further avoid outages. Establishing redundant power supply, however, is a topic on its own and not covered by this Master's Thesis.

6.2 Considerations

In this section we will make considerations on redundancy for the different parts of our test system. Please notice that we omit a solution using only one personal computer, since, in our opinion, this is the weakest part of our system, and numerous software parts (e.g. the bridging programs) cannot be installed/run twice in order to increase availability. We therefore aim at two PCs right away.

6.2.1 AC160 Side Devices

The AC160 controller and AF100 fieldbus provide a sophisticated redundancy concept. Thus, it is very common to insert two redundant processor modules as well as two communication interfaces into one rack. For other bus participants this configuration works transparently, so they do not care which PM/CI combination is currently in use and which is in stand-by mode. All communication interfaces already support line redundancy with two ports, this is also true for the CI527 PCI cards.

Not only controllers are included in the redundancy concept, but also OPC servers. That is, if two personal computers with CI527 and AC100 OPC Servers are connected to the same AF100 bus, both servers can provide data sent from the controller. Unfortunately, the redundancy concept is limited as this will not work in the opposite direction when using DSPs, which are vital for our application: Only one of the two servers is allowed writing to a specific variable on the controller. However, since we need to transmit a lot more values from AC160 to AC800M than vice versa, the OPC server redundancy concept helps us a great deal anyway: We do not have to define and send all these values twice (it is to mention here that defining and sending variables in AC160 is considerably less comfortable than in AC800M).

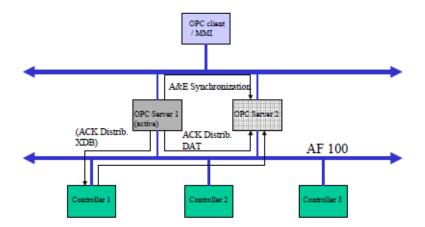


Figure 6.2: Redundant OPC server synchronization support [27]

In contrary, if we want redundancy for receiving variables, we have to define the receiving DSPs twice. The according redundancy brainpower must then be placed in the controller program code itself, that is, the code has to compare each received data pair and then decide how to proceed. In other words, the redundancy of the OPC servers/PCI cards is *not* transparent from the view of the AC160. Achieving

this transparency would only be possible by modifying the OPC server to support transparent redundancy when writing DSPs to the controller, or by the development and insertion of a specific hardware device between the AC160 and the PC.

6.2.2 AC800M Side Devices

Similar to the AC160, the AC800M supports processor module and communication interface redundancy for PROFIBUS as shown in Figure 6.3. In order to increase the availability of the CEX bus, also the redundancy link module BC810 is inserted [10].

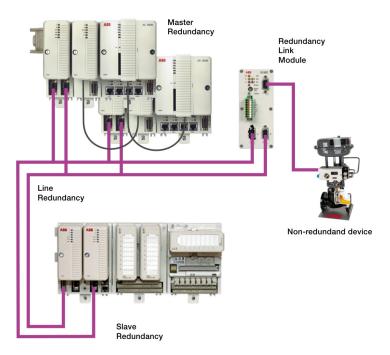


Figure 6.3: PROFIBUS redundancy concepts [37]

In contrary, the Woodhead PROFIBUS card and OPC server does not support slave redundancy and therefore limits the concept. It is to add, that *no single* PCI card/OPC server combination product that supports slave redundancy could be found on the market. However, Woodhead offers an API to control their PROFIBUS multi-slave card. That is, it would be possible to program an own redundancy logic for the PCI card. The OPC server, which has to support redundancy as well, needs to be adjusted or reprogrammed. Implementing this redundancy brainpower in the personal computer, one of the two cards has to be the active card while the other is in standby mode to appear transparently to the AC800M bus master.

Since the SST-PBMS-PCI card neither does support line redundancy, we would have to insert two cards in one computer and implement the line redundancy logic as computer software. As a far less laborious alternative, the RLM01 introduced earlier can be inserted.

The automation company Comsoft offers a hardware device called PRS to implement master redundancy in PROFIBUS DP systems [38]. It is intended to connect two masters with the same address to one PROFIBUS. If one master fails, the device, which contains a galvanic switch, changes to the other master. The switchover can also be initiated from outside via PROFIBUS itself or an additional Ethernet connection. The device is intended for implementing master redundancy, which is already done in our case. But, according to Comsoft, it is possible to "misuse" the device for implementing slave redundancy and therefore replacing a software solution [39]. However, logic to force a switchover would be needed anyway in order for the system to work properly.



Figure 6.4: Comsoft's PROFIBUS Redundancy Switch PRS

6.2.3 OPC Communication

It is of course necessary that the OPC servers for both bus interface cards are running, as well as the bridging software, e.g. Kepware LinkMaster. In addition it would be possible to add specific redundancy management software on the computers and interconnect them via (redundant) network and DCOM to the servers located on the second computer. That would theoretically allow the setup of a completely redundant system, coping with the outage of two different OPC servers or PCI cards at the same time, for example. Figure 6.5 shows the setup of such a completely redundant system. The black lines stand for OPC communication while the three dotted arrows symbolize the redundancy communication (RedCom) between two components.

There are different redundancy manager programs (RedMgr) available on the market, for example Matrikon's Redundancy Broker [40] or Kepware's RedundancyMaster [41]. They act as a server on one side while acting as a client on the other side, able to connect to redundant servers. The decision which server has to be considered the active one depends on the quality of the OPC connection and/or can

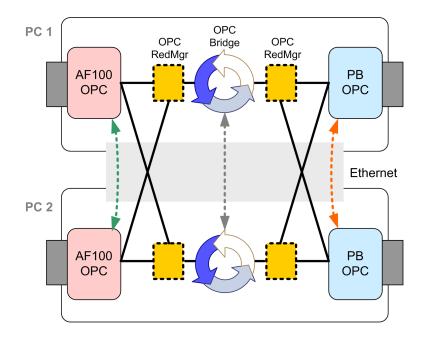


Figure 6.5: Diagram showing complete computer redundancy

be configured with additional checks of values. The software therefore also allows a restricted view on the transported data. However, these programs also insert an additional delay in communication, and from the computers point of view they are a weakness themselves since most likely all OPC connections are canceled if the program fails. Therefore it would also be possible to use different products at the same time, increasing the chance that one product keeps running when the other fails. However, a very interesting alternative is the software OPCFailover of Jemmac Software. While its function is the same as the classical redundancy managers, it does not insert a further relay, which also means no delay and no single point of failure. This is reached by just managing the OPC connection, always pointing to an active real OPC server, instead of acting as a proxy. As a drawback of this solution it is to mention that it does not support the periodical check of OPC values yet [42].

However, the introduction of such redundancy managers only makes sense if the OPC servers are implemented transparently redundant, meaning that only one sever is active at the time and without the need to send/receive some variables twice. This is not the case at the moment. The imaginary dashed green and red RedCom links in Figure 6.5 would have to be implemented first as discussed in the previous sections. Otherwise the OPC bridging software in PC1 will for example not match the variables defined in AF100 OPC Server of PC2. Dependent on the brainpower of both OPC server logics, maybe also the OPC bridges are not allowed to run concurrently and need an own (gray) redundancy communication. Matrikon's OPC Data Manager supports such a layout. Subsuming, there is one big advantage in implementing redundancy for the OPC connections in our system: It allows the simultaneous outage of several unequal parts of both computers. As a further consequence of the presumed transparent redundancy of the OPC servers, no variables have to be defined and sent twice in the controllers, there for saving transmission time. Unfortunately, there is also an considerable number of disadvantages:

- The support of transparent redundancy for both OPC server pairs is missing and has to be implemented first in order to establish a clean solution for OPC connection redundancy. This work is a huge programming task not to be underestimated.
- When using classical redundancy managers, also an extra delay and another OPC connection weakness is inserted. When using OPCFailover as redundancy manager instead, no periodical check of data can be implemented since this feature is missing.
- A save and possibly redundant network connection between the computers is needed.
- Only ODM supports a redundant setup if a redundancy link for the bridging software is needed. But our tests showed that ODM works slower than LinkMaster.
- It is possible that (e.g. due to a small temporary error) the current connection runs over the Ethernet network and in some cases even back again. This could make the connection slow.
- DCOM is considered as a weak part. We have to accept that or alternatively install even more software, namely a so-called OPC tunneler.
- If D/COM or other basic services fail, most likely all OPC communication on the computer will be canceled, no matter how redundancy is implemented.
- The complexity of the system makes it expensive as well as difficult to setup and observe.

6.3 Proposal

Unfortunately, neither the AC160 side of the system nor the AC800M side can be implemented with transparent redundancy on bus level. In both cases it is the PCI card/OPC server combination that lacks of this functionality. It is theoretically possible to add this functionality, but as there are numerous drawbacks in doing so and since there are other approaches, it is not recommended to do so.

We recommend using two personal computers as "single lines" without communication between the redundant counterparts. This saves implementation and programming time as well as communication delay due to relays. With this configuration, the system is not completely redundant, thus it is not allowed that more than one computer is affected by errors. However, in this approach we look at the computers as one functional device, and it is in fact very likely that if one application fails, the whole system is inconsistent and communication has to be switched to the second anyway.

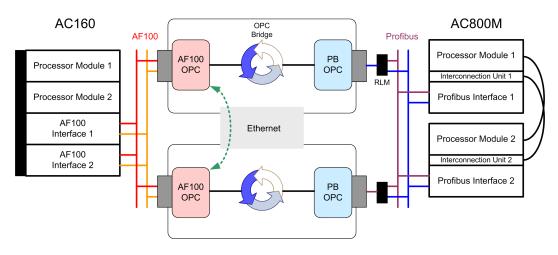


Figure 6.6: Proposal for redundancy

Using this configuration, the redundancy brainpower is located in the application programs of the controller. The two personal computer connections run at the same time. Each controller side puts their sending data on both channels and the controller on the other end has to select which information should be processed. This inspection on application level has also the advantage, that almost any kind of communication problem is detected. On the AC160 this kind of selection has to be implemented for every receiving signal. However, there are just 56 signals to proceed. The hundreds of sending signals do not matter since due to the AC160 redundancy concept they are distributed to both OPC servers anyway. On the AC800M the selection has to be implemented for every receiving signal, too. Since AC800M allows object oriented programming it is possible to program a generic code to be used for every signal. In contrary to AC160, also the sending signals have to be sent two times (to each PROFIBUS-card). The voting of receiving signals in both controller types can happen using different methods. One could for example check the quality of one or more specific variable/s for each channel and thus determine which of the two channels is used for incoming values. Alternatively it would be possible to always compare the redundant variables of both channels one after one and generating a resulting value using some voting algorithm.

If there is a physically long distance between the personal computer location and the AC800M, we recommend the use of the RLM01 in order to establish line redundancy. To keep things simple and since the device causes a small delay, it is reasonable to do without line redundancy on PROFIBUS if the devices are placed very close or even next to each other.

6.4 Summary

We realized that the subsequent implementation of redundancy in a mostly notredundant system is complex, difficult and also expensive. Furthermore, it slows down the connection due to introduced redundancy components and/or the twofold transmission of data. However, we proposed a solution that minimizes this drawbacks to acceptable size and can be implemented in a realistic time-frame. Nonetheless it is recommended to introduce redundancy only if it is really needed.

Finally we would like to mention that our considerations also apply when establishing a connection with MMS to AC800M instead of using PROFIBUS. The AC800M OPC Server does not support redundancy for writing variables, neither. However, similar to AC100 OPC Server, all (sending) variables can be accessed on both servers and the sending variables therefore need not to be sent twice.

Chapter 7

Outlook and Conclusion

This chapter will give an outlook on possible future work and summarize our thesis.

7.1 Contributions

In this thesis we make the following contributions:

- Evaluation and implementation of hard- and software to establish the connection.
- Quantitative results regarding the performance of a fast controller-to-controller connection via OPC.
- Evaluation of multiple configurations and their impact on performance and availability.
- Considerations and proposal regarding redundancy to improve availability and reliability.
- Identification, comparison and rating of different methods to connect controllers.

7.2 Outlook

This section first covers some hazards to be considered. The following part is divided into suggestions for improvements of the existing infrastructure on one hand, and in proposals for alternative approaches on the other hand. While some of the suggestions could be implemented in a comparatively short time, others require further investigation and long implementation times. We did not focus on the Advant Controllers since they are in the maintenance phase of their life-cycle. This means that that no developments will be done for these products.

7.2.1 Hazards

- In the PROFIBUS configuration tools as well as in the AC800M there are only limited capabilities for mappings of variables. That is, 32-bit values are either mapped to four 8-bit bytes or two 16-bit words. As this already happened for 32-bit floats and 32-bit integers (packed booleans!) in our configuration, no 32-bit mappings are left anymore. If we would like to process integers in addition, they are limited and mapped to 16-bit for example, or alternatively they would have to be split in booleans first (due to the transparent packaging in AC800M).
- While this was not considered a problem in our thesis, it could be one in a different configuration: If besides CI854A additional communication interfaces are connected to PM864A, the CEX bus could be overloaded and slow down communication.
- According to the specification, the length of a PROFIBUS segment at 12 Mbit/s baud rate can reach up to 100 meters. However, since a configuration at this maximum baud rate is considered to be rather sensitive, we propose to keep the distance as short as possible.
- Since our system implies various interfaces and open standards, it can be manipulated quite easily by anyone with sufficient know-how. In an operational environment it would thus make sense to limit access to the personal computer at least, securing login and network.

7.2.2 Improvements

General improvements

- To improve determinism one could replace the current Windows Server 2003 R2 operating system by a real-time version of Windows, for example with real-time extensions, Windows XP Embedded or Windows CE.
- As AC100 OPC Server offers stand-alone functionality, the system can be reduced to its vital parts, omitting all other tools: Operating system, OPC servers, bridging software.
- Signals could be split in different priority groups. Less important signals are processed slower on AF100 as well as in the OPC bridging software, therefore leaving more resources for the important signals which then can be processed faster.
- A combination of PROFIBUS and MMS would be possible, processing fast signals through PROFIBUS and slow ones via MMS, for example.
- A general reduction of signals would also improve performance, of course.

- Implementing the OPC Data eXchange standard in all OPC servers would supersede a bridging software. Since this standard was not popular until now, it is probably not an option for most manufacturers.
- With the new generation of OPC, Unified Architecture, a lot of improvements are implemented, for example COM-independency, redundancy support and security. We hope, therefore, that vendors will provide OPC UA servers for their devices as soon as possible.

PROFIBUS improvements

- PROFIBUS parameters were set automatically by AC800M. It should be possible to optimize these parameters in order to decrease the cycle time. However, the impact on the resulting RTT will be rather small.
- From our point of view, it should be possible for the manufacturers to implement the ability of swapping words/bytes either in CI854A or STT-PBMS-PCI with reasonable expenditure. This would save the unpleasant word swaps in AC800M program code and we can therefore only hope that this feature will be implemented in a future firmware release.
- Further development of CI854A could include the implementation of (multi-) slave functionality, thus offer new possibilities for PROFIBUS connection, especially the direct connection to PROFIBUS masters without the need for DP/DP-couplers. For example, there are master PCI cards supporting redundancy [43].

MMS improvements

- While Gigabit-Ethernet is already an accepted standard in office IT, AC800M provides only a 10 Mbit/s interface. Increasing this data rate would boost the performance of this connection. As a precondition, all involved parts need to support this higher rate. This applies to the processor module CPU in particular, since it is responsible for this communication.
- As a further step, the update rate of AC800M OPC Server should be implemented more flexible, such that an update rate down to 1 ms is possible if the rest of the communication permits to do so.

7.2.3 Alternatives

• Instead of using OPC as communication interface, the relay could be realized with own "glue logic" software. Woodhead offers an application programming interface with its card, for example. However, the implementation would be time-intensive and far less portable/flexible. Of course, also different operating systems like Linux could be taken into account, yet the card drivers must be written by oneself as far as we can tell.

- Even further goes the idea to implement the relay not on a personal computer, but on a customized hardware device containing a microcontroller/FPGA. This would result in so-called bus coupler device, which is commonly used for other bus types [44]. However, this approach is of course development intensive and very specific.
- With PM644 there exists a processor module for AC160 which provides a PROFIBUS interface [45]. This module's performance is too weak for the intended task, though. Nonetheless, it could be "misused" as communication interface, sending/receiving all data to/from the real processor modules PM665. Since it can only act as PROFIBUS DP master, just like CI854A, they need multiple DP/DP-couplers¹ as relaying devices. The lack of redundancy support could be compensated by the same measures as already discussed in the preceding chapter [46].
- Last but not least it would be possible to replace PROFIBUS and MMS by an alternative communication technology. Due to the restricted selection of communication interfaces and limited capabilities of (field) buses this does not make much sense today. But as *Industrial Ethernet* solutions like PROFINET [47] will become more and more important, this will be changing whithin the next years.

7.3 Conclusion

Results show that with the accurate components it is generally feasible to establish a fast (enough) controller-to-controller connection via OPC, meeting the requirements in average. With the Woodhead PROFIBUS solution, which showed best results, this can be reached using only one PCI card for each bus. However, the long-time test showed that there are limitations concerning the real-time requirements, since the maximum RTT we measured was almost three seconds. As long as this is the case, this approach cannot be used for critical applications. Moreover, there are smaller problems like the word swap that have to be considered. We are sure that this unwanted appearances can be eliminated by implementing some of the proposals of the preceding section. Nevertheless, the system can be used for in-house purposes, e.g. to test certain software parts before another approach is available, or even for non-critical operational applications like visualization tasks. Also the engineering effort to realize a PROFIBUS connection is acceptable, especially since variables can be processed transparently in AC800M. The use of tools like Excel and Bulk Data Manager is a big help in this context.

It turned out to be a big challenge to define a good redundancy concept for existing

 $^{^1 11}$ devices for the current quantity. Unfortunately, a multi-slave DP/DP-coupler could not be found until now.

infrastructure. Implementing complete redundancy even inside the personal computers seems to be in no relation to the effort. The proposed solution, treating the computers as single "lines" and moving redundancy brainpower to the controllers, is a realistic and feasible approach. However, it is engineering and programming effort anyway and redundancy can slow down the connection. It is therefore recommended to do without redundancy as long as such a system is only used for not process relevant in-house purposes.

Since OPC DA is based on Microsoft's COM, a proprietary solution, one will always be restricted in various directions. Along with many other improvements, OPC UA will get rid of this drawback and open up new perspectives to connect controllers. It is therefore important for ABB to build up know-how for OPC UA and implement these specifications in their own servers. An alternative, more specific and less general but also promising approach is the use of PM644 with AC160, which in our opinion should be tested out if the need for such an application persists. We believe that the testing of such a configuration can be done in a comparatively short period. Last but not least we would like to remark that importance of fieldbuses will decrease due to the entry of Industrial Ethernet solutions. We therefore recommend keeping an eye on this development.

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Appendix A Abbreviations

ADS	Automation Device Specification, a Beckhoff proprietary protocol
AE	Alarm and Events, an OPC specification
AF100	Advant Fieldbus 100
AMPL	Advant Master Programming Language
API	Application Programming Interface
BCB	Bus Configuration Builder
BIOB	Backplane Input/Output Bus
CBA	Control Builder A
CBM	Control Builder M
CDP	Cyclic Data Packet on AF100
CEX	Communication Extension Bus
CI	Communication Interface
CLC	Closed Loop Control
COM	Component Object Model, a Microsoft technology
CPU	Central Processing Unit
CSV	Comma Separated Values
DA	Data Access, an OPC specification
DAT	Database Element in AC160, e.g. Integer or Real
DCOM	Distributed COM
DCS	Distributed Control System
DEPP	Daten-Erfassungs-Programm P13
DP	Decentralized Peripherals, a PROFIBUS specification
DSP	DataSet Peripheral
EIA	Electronic Industries Alliance
FCB	Function Chart Builder
FMS	Field Message Specification, a PROFIBUS specification

FPGA	Field Programmable Gate Array
GUI	Graphical User Interface
IP	Internet Protocol
IT	Information Technology
MDT	Mean Down Time
MMI	Man Machine Interface
MMS	Manufacturing Message Specification
MTBF	Mean Time Between Failure
MVB	Multipurpose Vehicle Bus
ODM	OPC Data Manager, Matrikon's OPC bridging software
OLC	Open Loop Control
OLE	Object Linking and Embedding, a Microsoft technology
OPC	No Meaning, anciently OLE for Process Control
PA	Process Automation, a PROFIBUS specification
PB	PROFIBUS
PC	Personal Computer, not used as "process control" in our thesis!
PCI	Peripheral Component Interconnect
PLC	Programmable Logic Controller
PM	Processor Module
RTT	Round-Trip Time
SCADA	Supervisory Control And Data Acquisition
SDP	Service Data Protocol
TCP	Transmission Control Protocol
UA	Unified Architecture, OPC UA is a new generation of OPC

Appendix B

Version Numbers

Industrial IT 800xA	5.0
Control Builder M	5.0.0/0
PM864A firmware	5.0.11.61 (BasicHwLIb 5.0-0)
CI854A firmware	5.50-0
	(CI854PROFIBUSHwLib 1.0-19)
PM665 BASE Software	2.2/2
AMPL PC and DB Element Libraries	3.2/1
Control Builder A	
	1.2/1
Application Builder Function Chart Builder	4.1/0
	6.2/1
Bus Configuration Builder	3.0/1
FC3102 PCI card hardware	4
FC3102 PCI card firmware	2.58
TwinCat software	2.10.0
SST-PBMS-PCI firmware	1.3
SST-PBMS-PCI multi-slave module	1.24
SST Profibus Configuration	3.8
Multi Slave OPC Configuration Utility	1.5
Matrikon OPC Data Manager	5.8.0.0
Kepware LinkMaster	2.20.108
-	
MatrikonOPC Explorer	3.3.2.0
Microsoft Office	11
Process Explorer	10.21
Synergy	1.3.1
Irfanview	3.99

Appendix C

Task Description

Diplomarbeit ABB Schweiz AG

Power Systems / Power Generation

Industrielle Netzwerke Leistung – Verfügbarkeit – Zuverlässigkeit

Es sollen die bestehenden Strukturen von industriellen Netzwerken und Bussystemen für die Kraftwerksautomation basierend auf dem ABB Leitsystem 800xA analysiert werden. Dabei sollen verschiedene Fragestellungen in den Bereichen Verfügbarkeit, Leistung und Sicherheit bearbeitet werden.

Das ABB Control-System welches zur Steuerung von Hochleistungs-Gasturbinen eingesetzt wird, soll mit den neuesten Technologien punkto offener Schnittstellen, Anbindung an TCP/IP Netzwerke und Unterstützung von Device Management Tools etc. ausgerüstet werden. Hierzu soll der bestehende "Kopf-Rechner" Advant Controller 450 durch einen Controller der neuesten Generation, einen AC800 M Controller ersetzt werden.

Die Aufgabe besteht darin, den AC800 M Controller an einen Advant Controller 160 über das bestehende Interface AF100 OPC Server anzukoppeln. Es ist eine Lösung zu erarbeiten, welche zyklischen Datenaustausch des AC800 M Controllers mit dem AC160 Controller unter Einhaltung des geforderten Datendurchsatzes, der Verfügbarkeit und der Zuverlässigkeit ermöglicht. Dabei sind die zurzeit verfügbaren Produkte und Technologien der Plattform zu nutzen.

Die Bearbeitung der oben aufgeführten Fragestellungen beinhaltet sowohl einen theoretischen als auch einen praktischen Teil. Dabei sollen Daten für die Leistung, Verfügbarkeit und Zuverlässigkeit der erarbeiteten Lösung theoretisch hergeleitet, sowie ein System aufgebaut werden. Am aufgebauten System soll anhand von durchgeführten Tests Verbesserungsvorschläge resp. Änderungen an der System-Architektur evaluiert werden können.

Die technischen Spezifikationen des Control-Systems sind soweit möglich einzuhalten resp. Abweichungen entsprechend zu dokumentieren.

Eine optionale Ausweitung der Aufgabenstellung besteht in folgender Leistung. Das System 800xA basiert seit der Einführung vor mehreren Jahren auf OPC. In der Vergangenheit wurden solche Systeme immer in einem eigenen Netzwerk und eigener Windows Domäne realisiert. Heutige Kundenanforderungen zielen immer häufiger darauf ab, dass sich OPC Server und OPC Client in unterschiedlichen Netzwerken resp. Domänen befinden. Im Rahmen der Diplomarbeit soll geklärt werden, welche Einbussen in der Übermittlungsleistung in Kauf genommen werden müssen, wenn Server und Client in unterschiedlichen Domänen liegen resp. welche Massnahmen ergriffen werden können, damit die Leistungseinbussen möglichst gering bleiben.



Appendix D

Variables Overview

Test 1 (Beckhoff, Woodhead, MMS)

	Α	C160 to AC	800M	AC8	00M to AC	160	
	Status	Values	Total	Command	Values	Total	Grand
	Bits			Bits			Total
Count	992	10	1002	310	10	320	1322
Bits	1024	320	1344	320	320	640	1984
Bytes	128	40	168	40	40	80	248
Modules	32	10	42	10	10	20	62

10 DSPs with cycle time 8 ms: AF100 bus load 27%

Test 2 (Beckhoff, Woodhead, MMS)

	Α	C160 to AC	800M	AC8	00M to AC	160	
	Status Bits	Values	Total	Command Bits	Values	Total	Grand Total
Count	1984	20	2004	620	20	640	2644
Bits	2048	640	2688	640	640	1280	3968
Bytes	256	80	336	80	80	160	496
Modules	64	20	84	20	20	40	124

17 DSPs with cycle time 8 ms: Bus load 50%

Test 3 (Beckhoff, Woodhead, MMS)

	AC	C160 to AC8	300M	AC8	00M to AC	160	
	Status Bits	Values	Total	Command Bits	Values	Total	Grand Total
Count	2852	30	2882	930	30	960	3842
Bits	2944	960	3904	960	960	1920	5824
Bytes	368	120	<u>488</u>	120	120	240	728
Modules	92	30	122	30	30	60	182

24 DSPs with cycle time 8ms: Bus load 75%

Test 4 (Woodhead, MMS)

	AC	160 to AC8	00M	AC8	00M to AC1	60	
	Status Bits	Values	Total	Command Bits	Values	Total	Grand Total
Count	2170	600	2770	744	32	776	3546
Bits	2240	19200	21440	768	1024	1792	23232
Bytes	280	2400	2680	96	128	224	2904
Modules	70	600	670	24	32	56	726

91 DSPs with cycle time 64ms: Bus load 39%

Test 5 (Woodhead, MMS)

	AC	C160 to AC	800M	AC8	00M to AC	160	
	Status Bits	Values	Total	Command Bits	Values	Total	Grand Total
Count	4340	1200	5540	1488	64	1552	7092
Bits	4480	38400	42880	1536	1048	3584	46464
Bytes	560	4800	5360	192	256	448	5808
Modules	140	1200	1340	48	64	112	1452

182 DSPs with cycle time 128ms: Bus load 38%

Appendix E

Results Overview

						Rour	d Trip	Times	Round Trip Times in Milliseconds	lisecor	spu									
Using Kepware LinkMaster / Configuration C / 500 samples	vare Link	Master / C	onfigura	ation C /	500 sam		1/Test C) corresp	onds to	Test 1 w	ithout ch	anges / I	Nore rest	each / Test 0 corresponds to Test 1 without changes / More results and can be found in detailed logs	an be fo	und in de	etailed lo	SD		
		$\left \right $													$\left \right $		\parallel			
		Test 1	-			Test 2	5			Test 3	t 3			Test 4	4			Test 5	5	
	Min	Max M	Median	S.D.	Min	Max M	Median	S.D.	Min	Max	Median	S.D.	Min	Max N	Median	S.D.	Min	Max N	Median	S.D.
Complete System																				
Beckhoff without changes	30	100	60	9.6	40	06	60	10.3	40	100	70	11.3								
Beckhoff with changes	30	100	170	54.8	330	400	350	9.1	480	730	510	16.4								
Woodhead without changes	40	06	80	13.7	40	06	80	13.9	40	06	8	14.3	40	120	80	12.8	80	400	180	34.8
Woodhead with changes	40	06	80	14.5	50	120	80	12.3	40	100	8	8.8	40	140	06	17.1	110	1070	320	102.4
MMS without changes	100	260	170	35.1	100	240	170	31.6	110	250	170	31.3	100	230	120	29.2	110	550	230	65
MMS with changes	100	670	180	56.3	100	490	190	58.1	110	620	180	51.9	06	310	181	43.0	110	066	430	131.4
Shorted Controllers																				
AF160 without changes	14	80	40	10.6	13	80	32	10.0	17	104	40	11.2								
AC160 with changes	14	64	40	8.8	13	56	40	9.8	33	72	40	7.7								
Beckhoff without changes	20	80	30	7.8	10	60	30	8.0	20	60	ß	8.1								
Beckhoff with changes	60	170	80	8.4	130	250	160	10.8	210	350	230	15.3			-				-	
Woodhead without changes	40	70	50	4.6	40	60	50	5.7	40	60	50	4.5								
Woodhead with changes	40	60	50	4.6	50	80	50	4.1	40	80	50	4.0								
MMS without changes	40	390	120	34.4	40	240	120	31.6	40	190	120	31.8								
MMS with changes	80	310	120	27.8	50	190	120	35.7	100	370	120	22.2		-				-		
Additional																				
MMS time groups w/o changes					100	250	170	31.8	100	360	170	35.6								
MMS time groups with changes					100	560	190	53.8	110	490	230	50.5								

08.06.2007

Appendix F CD-ROM